Asymptotická složitost:

$f \in \Omega(g) \rightarrow$ algoritmus lze zdola odhadnout funkcí g

$f \in O(g) \rightarrow$ algoritmus lze shora odhadnout funkcí g

nulový prvek, většinou uvažujeme 0

$$f \in O(g) \Longleftrightarrow \exists c > 0, \exists n_0, \forall n > n_0: 0 \leq f(n) \leq c \cdot g(n)$$
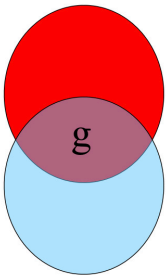
$$f \in \Theta(g) \Longleftrightarrow f \in O(g) \wedge f \in \Omega(g)$$

Théta   Oh   Omega

- The Sets big oh O(g), big theta Θ(g), big omega Ω(g)

Ω(g): functions that grow **at least as fast** as g

Θ(g): functions that grow **at the same rate** as g

O(g): functions that grow **no faster** than g

$$f \in \Omega(g) \Longleftrightarrow \exists c > 0, \exists n_0 \in \mathbb{N}: \forall n > n_0: 0 < c \cdot g(n) \leq f(n)$$